State preparation and operator growth of SYK model on IBM quantum computer

arXiv: 2311.17991 (Phys.Rev. D 109, 105002) [with Asad, B. Sambasivam] arXiv: 2406.15545 [in review, with Jack Araz, F. Ringer, B. Sambasivam]

Tensor Network 2024 November 17

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Complexity



Misconception: QC can solve all problems

QCs.

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- 1703.00454
- ulletup).

It turns out that for majority of problems, quantum computers would do no better than classical computers. A major research direction is to understand which problems can be solved efficiently by

For example, we know that scattering in ϕ^4 can be solved efficiently by quantum computers arXiv:

Class of problems which are best suited for quantum advantage belong to complexity class BQP. For ex: Shor's algorithm. Also Grover's algorithm but not as nice as Shor's (only polynomial speed-





Outline of the talk

- Effectiveness of tensor networks MPS approximation and towards universal quantum
- Sachdev-Ye-Kitaev (SYK) model of holography
- Quantum gates and real-time evolution using quantum circuits
- SYK model with N = 6, 8 Majorana fermions on IBM quantum computers with error mitigation real-time dynamics and thermal state preparation
- Summary and future directions



Tensor networks

- i.e., they are less entangled.
- complicated network such as MERA.



The most efficient classical method of studying the properties of lower-dimensional systems is tensor networks. The idea is based on the fact that if the Hamiltonian is sufficiently local and gapped, then the relevant sector of the entire Hilbert space is a tiny region which satisfies area-law entanglement

In this case, the vector space of dimension d^N can be described by O(d χ^2) where χ is the bond dimension of the MPS. This prescription fails for gapless systems and has to be replaced by more







Classical to Quantum

- otherwise.
- \bullet
- efficiently than classical computers.



An important ingredient of numerical lattice formalism is Wick rotation. Can't use sampling methods

Tensor networks can help sometimes but they have their own limitations. We need new tools to understand real-time dynamics of interacting field theories or quantum many-body systems.

We require fundamentally <u>new</u> idea of computing [Manin, Benioff, Feynman et al., circa 1978] such that we can compute exp(-iHt) for a given H in terms of circuits exploiting features of QM more

Best-known tensor networks can't simulate!		
nceasing enangement		
works can simulate!		
Number of gubits		





Approaches to universal quantum computing

- most popular Superconducting and Trapped Ion.
- ۲ qubit approach. Error correction not that well-developed.
- Susskind Hamiltonian]



Qubit approach — Manipulate and utilize the two-state quantum system. More than dozen approaches. Two

Qumodes approach — Use photons (quantum harmonic oscillator), infinite-dimensional HS. Not as popular as

This talk will discuss the qubit approach, however, other approach might be better suited for bosonic d.o.f as explored for NLSM model (see 2310.12512). Now extending the "CV" approach to SU(2) gauge theory [Kogut-







Qubits vs. Qumodes

	CV
Basic element	Qumodes
Relevant operators	Quadrature operators \hat{x}, \hat{p} Mode operators $\hat{a}, \hat{a}^{\dagger}$
Common states	Coherent states $ lpha angle$ Squeezed states $ z angle$ Number states $ n angle$
Common gates	Rotation, Displacement, Squeezi Beamsplitter, Cubic Phase

Qubit

Qubits

Pauli operators $\hat{\sigma}_x, \hat{\sigma}_y, \hat{\sigma}_z$

Pauli eigenstates $\ket{0/1},\ket{\pm},\ket{\pm i}$

ing,

Phase Shift, Hadamard, CNOT, T Gate



Quick Recap - Unitary gates

$$|+\rangle = \frac{1}{\sqrt{2}} \left(|0\rangle + |1\rangle \right)$$



Operator	Gate(s)		Matrix
Pauli-X (X)	- X -		$\begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$
Pauli-Y (Y)	- Y -		$egin{bmatrix} 0 & -i \ i & 0 \end{bmatrix}$
Pauli-Z (Z)	$-\mathbf{Z}$		$\begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix}$
Hadamard (H)	$-\mathbf{H}$		$\frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1\\ 1 & -1 \end{bmatrix}$
Phase (S, P)	$-\mathbf{S}$		$\begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix}$
$\pi/8~(\mathrm{T})$	$-\mathbf{T}$		$egin{bmatrix} 1 & 0 \ 0 & e^{i\pi/4} \end{bmatrix}$
Controlled Not (CNOT, CX)			$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$
Controlled Z (CZ)			$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & -1 \end{bmatrix}$
SWAP		_*_ _*_	$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$
Toffoli (CCNOT, CCX, TOFF)			$\begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \end{bmatrix}$





<u>Questions?</u>

SYK model

$$H = \frac{(i)^{q/2}}{q!} \sum_{i,j,k,\cdots,q=1}^{N} J_{ijk\cdots q} \quad \chi_i \chi_j \chi_k \cdots \chi_q,$$

- ٠ $\overline{J_{...}} = 0, \ \overline{J_{...}^2} = \frac{q!J^2}{Nq-1}.$
- The fermions χ satisfy, $\chi_i \chi_i + \chi_i \chi_i = \delta_{ij}$. We will set J = 1. Note that it has units of energy and inverse time. •
- maximal Lyapunov exponent.



Model of N Majorana fermions with q-interaction terms with random coupling taken from a Gaussian distribution with

• In the limit of large number of fermions with $N \gg \beta J \gg 1$, the model has several interesting features such as



Mapping fermions to qubits

$$\chi_{2k-1} = \frac{1}{\sqrt{2}} \left(\prod_{j=1}^{k-1} Z_j \right) X_k \mathbb{I}^{\otimes (N-2k)/2}$$

- N Majorana fermions requires N/2 qubits. We use the standard Jordan-Wigner mapping to write χ in terms of Pauli matrices X, Y, Z, and Identity.
- . The SYK Hamiltonian is then written as sum of Pauli strings. The number of strings is $\binom{N}{q}$ a grows like $\sim N^q$. Simplest non-trivial case for is N = q with one term. We restrict to q = 4.

$$, \qquad \chi_{2k} = \frac{1}{\sqrt{2}} \left(\prod_{j=1}^{k-1} Z_j \right) Y_k \mathbb{I}^{\otimes (N-2k)/2}$$



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Simplest case: N=4

 $H = J_{1234} \chi_1 \chi_2 \chi_3 \chi_4$

- Hamiltonian. So, for this case we have $\exp(-iHt) = \exp(iJ_{1234}ZZt)$.
- This circuit is simple to construct and just needs 2 CNOTs and 1 rotation gate.



 $\chi_1 = XI, \, \chi_2 = YI, \, \chi_3 = ZX, \, \chi_4 = ZY$

 $H = J_{1234}(X\mathbb{I}) \cdot (Y\mathbb{I}) \cdot (ZX) \cdot (ZY) = -J_{1234}ZZ$

The goal of quantum computation is to construct a unitary operator corresponding to this





Circuit complexity

$$e^{-iHt} = \left(\prod_{j=1}^{m} e^{-iH_jt/r}\right)^r + \mathcal{O}\left(\sum_{j< k} \left| \left| \left[H_j, H_k\right] \right| \left| \frac{t^2}{r} \right),\right.\right.$$

Recall that N = 4 needed just 2 2q-gates for each Trotter step.

* Corollary of Zassenhaus formula i.e., $exp(t(X+Y)) = exp(tX) exp(tY) + O(t^2)$ (also known as dual of BCH formula).

Definition: How many 2q-gates do we need to simulate the SYK model?

Different approaches can be used to do the Hamiltonian simulation (aka time evolution). A popular method is Trotter method. It is based on Lie-Suzuki-Trotter product formula* (writing $H = \sum_{i=1}^{m} H_{j}$, $m \sim N^4$)

Depending on what error (ϵ) we desire in the time-evolution from the second term, we can compute the number of slices (r) we need to take. So, the complexity reduces to finding number of 2q-gates for each Trotter step.



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$$\mathscr{C} = \mathcal{O}(N^{10}t^2/\epsilon)$$
$$\mathscr{C} = \mathcal{O}(N^8t^2/\epsilon)$$
$$\mathscr{C} = \tilde{\mathcal{O}}(N^{7/2}t)$$

- basic idea of embedding H in a bigger vector space.
- lacksquarecomplexity to $\mathscr{C} = \mathcal{O}(N^5 t^2 / \epsilon)$ which we now discuss.



L. García-Álvarez et al., PRL 119, 040501 (2017) Susskind, Swingle et al., arXiv: 2008.02303 (2020) Babbush et al., Phys. Rev. A 99, 040301 (2019)

• The last one clearly is the most efficient, however, in the noisy-era implementing this is not feasible. It requires fault-tolerant quantum resources + ancillas since it is based on the

Using the Trotter methods, the best seems to be $\sim N^8$. In our paper, we improved the

Commuting terms

The costs can be simplified if we are little careful in splitting the SYK Hamiltonian.

can find diagonalising circuit for each cluster and then apply time-evolution operator.

This is in general a NP-hard problem but various approx. algorithms exits.

- The number of terms grows like $\sim N^4$, however, a large fraction of them commute with one another and can be collected together and then time-evolved more efficiently. We
- Finding optimal number of such clusters is a well-studied computer science problem.

Estimate based on general commutivity

\overline{N}	Pauli strings	Clusters	Two-qubit gates
4	1	1	2
6	15	5	30
8	70	6	110
10	210	23	498
12	495	57	1504
14	1001	92	3560
16	1820	116	6812
18	3060	175	11962
20	4845	246	19984

- evolution time t i.e., $\mathscr{P}_0 = |\langle \psi_0 | e^{-iHt} | \psi_0 \rangle|^2$. For initial state, we take $|0\rangle^{\otimes N/2}$.
- quantum circuit.
- straightforward. For N = 6, there are 30 2Q gates per step which we cannot show here.



Return probability

• A simple observable we can compute is the probability that we return to the same initial state after some

• For approximating the unitary, we use the first-order product formula and construct the corresponding

• For N = 4, we have a simple circuit of only two 2Q gates, so the entire circuit for return prob. is

$$R_z(-2J_{1234}\delta t)$$



IBM chip topology

combination of qubits we used for N = 8. This chip topology is 'heavy-hex'.



• We used the quantum computers available through IBM to simulate the SYK model. The topology of the processor is shown below. In practice, we need more gates than necessary. For example, we show a



<u>Return probability - IBM Results</u>





- coherence time of qubits.
- \bullet and extended in Rahman et al. arXiv:2205.09247

M3 is a matrix measurement mitigation (MMM or M3) technique that solves for corrected measurement probabilities using a dimensionality reduction step



DD (dynamical decoupling) — a series of strong fast pulses are applied on the system which on average increases the lifetime of qubits and delays decoherence (or effect of interactions with environment)

Error Mitigation

The results from the 127-qubit device (red) agrees slightly less than those with self-mitigation (green). The red points have been read from some fixed number of measurements/shots and postprocessed with mild mitigation including M3 to correct read-out errors and DD to increase

To get closer to the exact results, we found that an idea similar to CNOT only mitigation (known as self-mitigation) seems to help drastically. Basic idea introduced in Urbanek et al. arXiv: 2103.08591



SYK model - Bound on chaos

- SYK model famously saturated the Lyapunov exponent i.e., $\lambda = 2\pi T$ for $J/T \gg 1$ when N is large.
- OTOC := $\langle W(t)V(0)W(t)V(0)\rangle_{\beta}$ = Tr($\rho W(t)V(0)W(t)V(0)$) which characterizes quantum chaos.
- size ~billion. Hard for classical computers.

• One considers $C(t) = -\langle [W(t), V(0)] [W(t), V(0)] \rangle$ and the expansion of the commutator gives

• Suppose one starts at t = 0, and computes also the two-pt correlator given by $\langle W(t)W(0) \rangle$, the time scales at which the lower order correlators decay is called 'dissipation time'. After this time, the OTOC grows as $\exp(\lambda t)$ and saturates beyond t_{\star} known as scrambling time. Black holes are fastest scramblers!

These correlators have been computed up to N = 60 numerically i.e., H has ~million terms and matrix has



Out-of-time correlators (OTOC)

- So the goal is to compute $\langle W(t)V(0)W(t)V(0)\rangle_{\beta}$ on a quantum computer. Thermal correlators are currently not easy to compute due to limited resources. One simplification we can make is consider the $\beta \to 0$ limit of OTOC. This is not at all interesting for holography, but this is where we must start. Hence, the density matrix is just $\rho \propto \mathbb{I}$.
- The unusual time-ordering of OTOC is also hard for quantum computers which often mean carrying out forward and backward evolution. We use a protocol (next slide) which uses only forward evolution to compute OTOC on quantum hardware.



- ensemble (CUE).
- protocol works when W is traceless operator.

There are various protocols to measure OTOC on quantum computers, see Swingle 2202.07060 for review.

• We use the one proposed in 1807.09087 now known as 'randomised protocol' since it computes OTOC through statistical correlations of observables measured on random states generated from a given matrix

Infinite-temp OTOC is given by $Tr(W(t)V^{\dagger}W(t)V) \propto \langle W(t) \rangle_{\mu} \langle V^{\dagger}W(t)V \rangle_{\mu}$ where the average is over different random states $|\psi_{\mu}\rangle$ prepared by acting with random unitary on arbitrary state say $|0\rangle^{\otimes n}$. Note that this





Randomised Protocol



t-designs equivalent to first *t* moments of Haar group

We need two measurements (between which we compute the statistical correlation) and it is shown below. This is the global version of the protocol (since u has support over all qubits). There is also a local version of the protocol. Note that cost of decomposing arbitrary u increases exponentially, one can instead use unitary from a subset of Haar measure. They are called unitary t-designs* in literature.



 \checkmark



OC Results

lacksquaredid for return probability.



We used ibm_cusco and ibm_nazca to obtain the results show for N = 6. We took simplest operators where both W and V were taken to be single Pauli. We see good agreement without need to do self-mitigation like we

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Finite-temperature SYK model

- We considered OTOC measured over random states (maximally mixed) generated i.e., $\beta = 1/T = 0$. However, much of interesting Physics of the SYK happens in the region $\beta \gg 1$ and classical computations have argued that you need $\beta \sim 70$ to extract Lyapunov exponents close to the chaos bound.
- Finite-temperature OTOCs are difficult for quantum computers in general. No simple/general cost-effective protocol exists. To move towards this goal, we are studying the preparation of Gibbs (thermal) states on quantum computer for the SYK model.
- In addition to the thermal state (mixed) of the SYK model, one can also consider a purification of this known as thermo-field double state (TFD). TFD state is a pure state (up to unitary trans.) of some other system (for ex: coupled SYK model) and when we perform partial trace over either system, we recover thermal state on the other one.





- ground states on QC.
- parameters of a quantum circuit ansatz that minimizes the Hamiltonian expectation value.



VQE algorithm

Before we move to preparation of Gibbs state, let's us look at popular algorithm for preparing (approximate)

Hybrid classical/quantum algorithm to find the ground state problem of a given Hamiltonian by finding the

It primarily consists of three steps: 1) Prepare initial ansatz on QC i.e., $|\psi(\Theta)\rangle$, 2) Measure energy on QC and optimise the parameters Θ using classical optimisers and 3) Repeat until desired convergence is achieved.



<u>VQE for finite temperatures</u>



 \checkmark

arXiv: 2406.15545

Finite-temperature VQE methods are still an active farea of research. Many proposals exist. The cost function is no longer E but rather E - TS (free energy) which can be hard to compute on QC.

Selisko et al., 2208.07621





Finite-temperature SYK model



Results from PennyLane simulator





- We are entering an era where we can do small computations reliably on quantum computers. Exploring these toy models will hopefully reveal to us better algorithms/methods.
- It is important to note that if we can model the noise in these quantum devices, we can mitigate and get reasonable results!
- Preparation of thermal states and its purifications known as TFD states is still challenging and search for better ways to do this (by minimizing over gate costs and improving fidelity) is interesting research direction.

Summary



Resources and Data Statement

Published November 25, 2023 | Version v1

A model of quantum gravity on a noisy quantum gr

Asaduzzaman, Muhammad¹ (D; Jha, Raghav G.² (D; Sambasivam, Bharath³ (D

Additional resources for the arXiv article: https://arxiv.org/abs/2311.17991 including for details.

Files

OTOC_N6.zip	
OTOC_N6.zip	
► N=6	
H_N6_3.mtx	
H_N6_4.mtx	
H_N6_7.mtx	
C QC_N6_3.qasm	
C QC_N6_4.qasm	
C QC_N6_7.qasm	
ham_paulis_N6_3.txt	

Both classical and quantum code available at: https://github.com/rgjha/SYKquantumcomp

Computational notebook Computational Computa	Edit	
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	Versions	
×	Version v1 Nov 25, 2023 10.5281/zenodo.10202045	
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<u>Thank you</u>